

# A Simple Model of Smart Contracts in Agda

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- **Ethereum** = A **second-generation Blockchain** technology [7].
- Launched by Vitalik Buterin [4] in 2013.
- Main difference to Bitcoin is in the use of smart contracts:
  - Ethereum [9]:
    - ★ **Turing complete** language which includes loops;
    - ★ allows **calls to other contracts**;
    - ★ adds **cost of execution of instructions** (gas) to guarantee termination.
- Recently switch from **proof of work** to **proof of stake** [5], **solving the waste of energy problem**.

- **Smart contract** = **program** which is **automatically executed** when conditions in the blockchains are satisfied [8].
- Smart contracts are **immutable programs** [3].
- Smart contracts in the cryptocurrency Ethereum are usually written in the high-level language **Solidity** [6] which compiles into the low-level **Ethereum Virtual Machine (EVM)** [4].
- **World State Machine** with essentially immutable history.
- **Example applications:**
  - Tracing of goods (using we have an immutable database),
  - Electronic voting,
  - NFT (ownership of electronic items),
  - Investment funds (DAO).
- Because of **high monetary impact**, **immutability**, and **shortness of programs**, **prime candidate for verification**.

# Smart Contracts

- Blockchain is roughly speaking a **data base** which **determines for each address its current state** (amount of money, other data).
- In Ethereum **smart contracts = objects deployed to addresses**, with methods which can be called by
  - ordinary (externally owned) accounts,
  - other smart contracts.
- Toy example (Solidity):

```
1  pragma solidity ^0.8.17;
2
3  contract testLedger {
4      function f (int n) public pure returns (int){
5          return g(n);
6      }
7
8      function g (int n) public pure returns (int){
9          if (n > 0) {return f(n - 1);}
10         else      {return 0;}
11     }
12 }
```

- **Previous work:** Verification of Bitcoin smart contracts using **weakest preconditions** of Hoare logic [2, 1] in Agda.
- Goal of this and follow up papers is adaption to Solidity style smart contracts.
- **First Step Here:** develop model of Solidity-style smart contracts in Agda.
- More complex, because of use of **objects**.
- We cover execution of contracts including **calling of other contracts** and contracts having **multiple functions (methods)**.

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# Messages

- EVM allows calling functions with **serialised parameters**.
- Parameters represent elements of **arrays, maps, enumerations, integers**, etc.
- In our model, we abstract from this encoding by defining a **message data type**:

```
data Msg : Set where nat : (n : ℕ) → Msg
                    list : (l : List Msg) → Msg
```

- Arrays are represented as lists of messages.
- Maps are represented as lists of pairs (represented as lists) of messages.

# Programs (SmartContractExec)

```
data SmartContractExec : Set where
  return : Msg → SmartContractExec
  call   : SmartContractExecStep → SmartContractExec
  error  : ErrorMsg → SmartContractExec
```

- `SmartContractExec` consists of three constructors:
  - `return` = terminates execution and return its argument;
  - `call` = calls `SmartContractExecStep`  
then continues as defined by its continuation argument
  - `error` = raises an error.

# SmartContractExecStep

```
record SmartContractExecStep : Set where
  coinductive
  field calledAddress : Address
        calledFunction : FunctionName
        calledMsg      : Msg
        cont           : Msg → SmartContractExec
```

`calledAddress` = address of contract being called;  
`calledFunction` = function name called;  
`calledMsg` = argument of the function (a message);  
`cont` = continuation, depends on  
the result of executed function.

- `SmartContractExec` and `SmartContractExecStep` are defined coinductively, so loops and even non-terminating programs are allowed.

# Ledger and ExecutionStack

- A ledger determines for any address function name and msg argument the smart contract function to be executed:

$\text{Ledger} = \text{Address} \rightarrow \text{FunctionName} \rightarrow \text{Msg} \rightarrow \text{SmartContractExec}$

- **ExecutionStack** = stack of continuations
  - continuation are executed once the result of the execution above it has finished giving an element of **Msg**.

$\text{ExecutionStack} = \text{List} (\text{Msg} \rightarrow \text{SmartContractExec})$

# StateExecFun

- The state of execution is given by

```
record StateExecFun : Set where
  constructor stateEF
  field executionStack : ExecutionStack
       nextstep         : SmartContractExec
```

i.e. having two fields:

- `executionStack` is the current execution stack;
- `nextstep` is the current code to be executed.

# stepEF and stepEFntimes

- `stepEF`, is the one-step execution of a smart contract.
- `stepEFntimes`, which iterates it  $n$  times, corresponding to execution with a simple form of gas limit.

`stepEF` : `Ledger`  $\rightarrow$  `StateExecFun`  $\rightarrow$  `StateExecFun`

`stepEFntimes` : `Ledger`  $\rightarrow$  `StateExecFun`  $\rightarrow$   $\mathbb{N}$   $\rightarrow$  `StateExecFun`

$$\begin{aligned} \text{evaluateNonTerminating} : & \text{Ledger} \rightarrow \text{Address} \rightarrow \text{FunctionName} \\ & \rightarrow \text{Msg} \rightarrow \text{NatOrError} \end{aligned}$$

We can define as well a terminating version with additional parameter

$$gas : \mathbb{N}$$

which restricts evaluation to gas many steps.

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# Example of simple Solidity-style of smart contract in Agda

- Example recursively decrementing by 1 until 0:

```
testLedger : Ledger
```

```
testLedger 0 "f" (nat n)
```

```
  = call (smartContractExecStep 0 "g" (nat n) return)
```

```
testLedger 0 "g" (nat (suc n))
```

```
  = call (smartContractExecStep 0 "f" (nat n) return)
```

```
testLedger 0 "g" (nat 0)
```

```
  = return (nat 0)
```

```
testLedger ow ow' ow''
```

```
  = error (strErr " Error undefined")
```

- `evaluateNonTerminating testLedger 0 "f" (nat 5)`  
evaluates to `nat 0`

# Example in Solidity language

- Corresponding Solidity code:

```
1  pragma solidity ^0.8.17;
2
3  contract testLedger {
4      function f (int n) public pure returns (int){
5          return g(n);
6      }
7
8      function g (int n) public pure returns (int){
9          if (n > 0) {return f(n - 1);}
10         else      {return 0;}
11     }
12 }
```

# Example Run in Solidity

When applying "f" to 7 and "g" to 4 we obtain the following results:

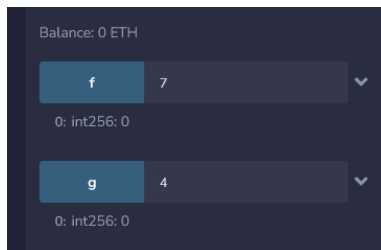


Figure: Result using Solidity language

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# Conclusion

- We developed a **simple model** of **Solidity-style smart contracts** in Agda.
- Dealt with **execution** and **calling of other contracts**.
- Not yet support of **gas cost, amount of money, transfer of money, state**.
- **Work in progress:**
  - Extend the simple model by the not yet supported items.
  - Develop an **interactive program in Agda** which allows to execute calls of functions in contracts with a corresponding ledger.
- **Future work:**
  - Adapt the verification of bitcoin using weakest preconditions [2] to verifying contracts in this model.

Thank you for listening.

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